

Hacking GeckOS

```
Init V1.0 booting
Start "fsdev ": ok!
Start "fsiec ": ok!
Start "shell b c:auto.bat ": ok!
Prepared restart!
Start "c:1sh -d c: ":
sh v0.1 21dec1997 (c) A. Fachat

> ok!
Prepared restart!

>uname
GeckOS/A65 2.0 6510 C64 1ib6502 0.6
>■
```

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World of Commodore
2019-12-07

Speaker Bio

- ✓ a.k.a. "Cenbe"
 - ✓ retired Java programmer/Linux sysadmin
 - ✓ collector of programming languages and operating systems for the Commodore 64:
- <https://www.lyonlabs.org/commodore/>

Happy 50th, Unix!



wait... "Unix" on a 6502?

ARE YOU

CRAZY?

wait... “Unix” on a 6502?

Let’s compare a Commodore 64 to the machines that early versions of Unix ran on:

	PDP-7	PDP-11	Commodore 64
memory:	16K	24K	64K

Processor speeds were comparable to that of a ‘64, but the architecture was very different.

Of course, there are the registers...

GeckOS history

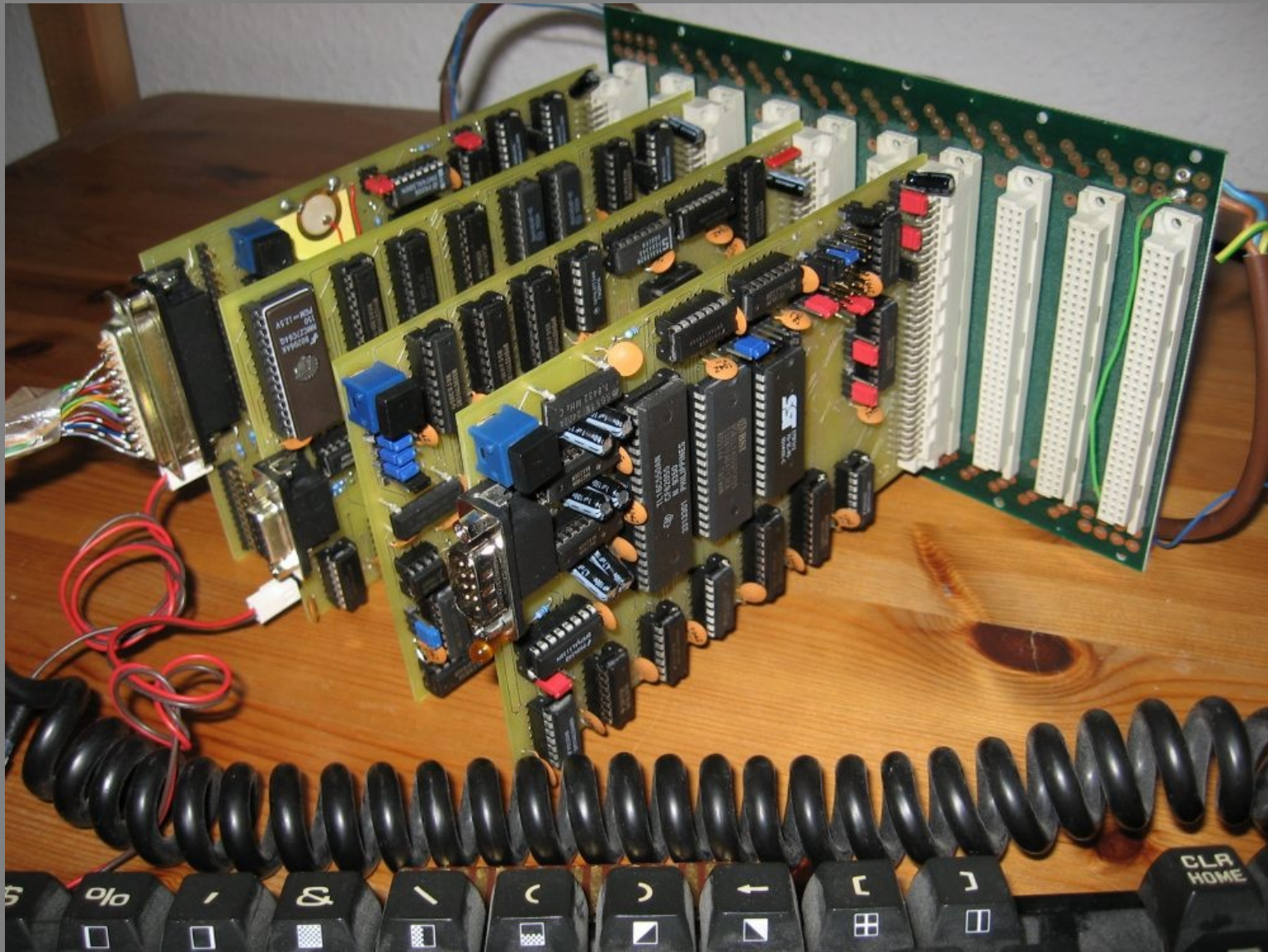
- ✓ Written by André Fachat, originally for a 6502 computer with an MMU that he built in 1989.
- ✓ He later ported it to run on the PET, 8296, and C64.
- ✓ Development is very active!
- ✓ Source is GPL, on GitHub (fachat/GeckOS-V2).



André Fachat

CS/A65

<http://www.6502.org/users/andre/csa/index.html>



GeckOS features

GeckOS is a Unix-like multitasking operating system for the 6502 CPU. It supports:

- ✓ task priorities, multi-threading
- ✓ virtual consoles
- ✓ signals, semaphores
- ✓ backgrounding and redirection
- ✓ piping
- ✓ environmental variables
- ✓ relocatable file format

DEMO

- ✓ shell (both), monitor
- ✓ forking (one program loads and runs another)
- ✓ backgrounding a program (“the Schema demo”)
- ✓ signals (sending messages between programs)
- ✓ semaphores (blocking on available resource)
- ✓ new `ps` and `kill` commands

This is a very recent build of GeckOS.

original info command (old shell)

```
SHELL V1.3  
(C) 1990-97 BY A.FACHAT
```

#	info									
PID	Name	Th	En	Pa	Me	Sm	SigA	In	Out	Err
00	1sh	01	00	FF	78	40	8A4A	FC	00	FC
0E		01	00	00	78	00	0000	FC	FC	FC
1C		01	00	00	78	00	0000	FC	FC	FC
2A		01	00	00	78	40	AFA1	01	02	02
38		01	00	00	78	4F	B4D9	03	00	00
46		00	00	00	78	00	0000	00	00	00
54		00	00	00	78	00	0000	00	00	00
62		00	00	00	78	00	0000	00	00	00
70		00	00	00	78	00	0000	00	00	00
7E		00	00	00	78	00	0000	00	00	00
8C		00	00	00	78	00	0000	00	00	00
9A		00	00	00	78	00	0000	00	00	00
00		00	00	00	00	00	0000	00	00	00
00		00	00	00	00	00	0000	00	00	00
00		00	00	00	00	00	0000	00	00	00

```
#
```

new ps command (lsh shell)

```
Init V1.0 booting
Start "fsdev ": ok!
Start "fsiec ": ok!
Start "shell b c:auto.bat ": ok!
Prepared restart!
Start "c:lsh -d c: ":
sh v0.1 21dec1997 (c) A. Fachat

> ok!
Prepared restart!

>ps
PID  Name      Exec  Th  Pa  Sw  SigA  In  OutErr
00  init      8721  01  FF  40  8A4A  FC  00  FC
0C  fsdev     8B09  01  00  00  0000  FC  FC  FC
18  fsiec     8EAD  01  00  00  0000  FC  FC  FC
24  shell     97E2  01  00  40  AF70  01  02  02
30  c:lsh     184C  01  00  4F  B4C1  03  00  00
3C  ps        237B  01  30  4F  B4C1  03  00  00

>■
```

GETINFO and the task table

- ✓ `info` (the old shell's "`ps`") calls the kernel `GETINFO` API, which reads the task table and returns information about all processes.
- ✓ It uses the program communication buffer (`PCBUF`, a.k.a. `SYSBUF`) to build a table (since programs should not have direct access to the task table).
- ✓ The task table did not originally have entries for either process name or exec address, although the `GETINFO` table has one for name.

adding process names (stdlib programs)

- ✓ For lib6502 programs, the name can be found in the LIBSAVE structure which is populated when a program is started. This structure is pointed to from the task table.

PROBLEM: the kernel shouldn't assume that programs are written using the standard library, and shouldn't access the LIBSAVE structure, as it is lib6502-specific.

adding process names (kernel programs)

- ✓ For init and the device drivers, it's possible to get the name by walking the ROM image headers in the same order that kernel startup does.

PROBLEM: this breaks if kernel initialization changes... it's also a filthy kludge!

Let the kernel do it!

SOLUTION: the kernel `FORK` routine takes process name and exec address as parameters; it should just save them in the task table.

PROBLEM: lib6502 programs pass the program name with a stream number in the first byte.

SOLUTION: change lib6502 to pass the stream number as a parameter to `FORK` (kernel passes this byte back in `.A` when the process starts).

adding exec address (kernel programs)

- ✓ The program headers in the kernel image contain the exec address, so it's an easy matter for the kernel to put it in the task table when starting one of these programs.

adding exec address (stdlib programs)

- ✓ The exec address is passed to FORK, and could be stored just before it passes control to the program.

PROBLEM: lib6502 programs set a start address of lib6502's `libfork` routine (which loads and relocates the program).

SOLUTION: provide a `SETINFO` API that would allow lib6502 to update the task table after `FORK` has been called.

A New Golden Age for GeckOS

Now that we can debug more easily, anything is possible:

- ✓ The Grand Unification of the Shells
- ✓ better support for CMD HD, μ IEC, 1541 Ultimate (partitions, subdirectories, disk images...)
- ✓ 1541 Ultimate networking
- ✓ native speeder in the filesystem?
- ✓ *your project here*

resources

- ✓ GeckOS (source, tools, docs, disk images):
<http://www.6502.org/users/andre/osa/index.html>
- ✓ GeckOS source on GitHub:
[fachat/GeckOS-V2](#)
- ✓ online HTML documentation:
<https://www.lyonlabs.org/commodore/onrequest/GeckOS-docs/index.html>
- ✓ Cenbe's Commentary on GeckOS:
<https://www.lyonlabs.org/commodore/onrequest/geckos-analysis.html>

QUESTIONS